RADIO NETWORK, RELAY NODE, CORE NODE, RELAY TRANSMISSION METHOD USED IN THE SAME AND PROGRAM THEREOF

BACKGROUND OF THE INVENTION

Field of the Invention

The present invention relates generally to a radio network, a relay node, a core node, a relay transmission method used in the same and a program thereof. More particularly, the invention relates to a relay route setting method and a relay transmission method in a cellular system, in which a plurality of nodes are connected by a radio network.

Description of the Related Art

A cell structure in the conventional cellular system is illustrated in Fig. 20. In Fig. 20, the reference numeral 401 denotes a cell, 402 denotes a base station (node). As shown in Fig. 20, a service area is constructed by arranging a plurality of cells as shown in Fig. 20.

Each node are connected by a wired backbone network 404 through wire lines 403 for transmitting service signal, such as voice, data and so forth and various control signals. It should be noted that the node and the wired backbone network may be connected by hierarchically providing line concentration station therebetween.

A terminal station 405 communicates with a node 402 for

performing transmission and reception of various signals to be transmitted by a wired backbone network 404 through wire lines 403. In the wired backbone network, not only the radio base station (node) but also a server equipment performing management of position information of the terminal station 405, billing process and so forth is provided.

In order to adapt for increasing of subscribers of cellular system, such as cellular telephone, subscriber fixed wireless access and so forth, radius of each cell is reduced to reduce process load in each node. When the system is constructed with such micro cells, quite large number of nodes are arranged in order to certainly establish the service area.

On the other hand, when a high density data transmission method of multi-value modulation or the like is applied for adapting to a high speed data transmission, service area to be covered by one node inherently becomes narrow for certainly quality. Even in such case, quite large number of nodes have to be arranged for certainly establishing service area.

Furthermore, the conventional cellular system has been mainly designed in a quasi-microwave band and microwave band, due to tightness of frequencies, it has been desired to establish the cellular system using submillimeter wave or millimeter wave. When frequency becomes higher, diffraction effect of the radio wave is reduced to make straightening characteristics

significant, and it causes difficulty in non line-of-sight communication to inherently narrow area of service in each node. Namely, even in such case, the service area has to be certainly established with micro cells to install quite large number of nodes.

When system is established with large number of micro cells, it becomes essential to establish wired network for connecting the node group to the backbone network. However, for connection between quite large number of geographically concentrated nodes and the backbone network, wire line networks have to be extended to every places to inherently increase cost of the overall system. Therefore, there is a method to establish connection between nodes with wireless communication for relay transmission to expand service area.

Setting of the relaying dominates a capacity because capacity of the cellular system is restricted by interference, and tolerance to interference depends upon the setting of the relaying. In route method that makes number of relay nodes minimum, so-called minimum hop routing method, throughput or system capacity cannot be always maximum as viewed overall relaying route for lacking of reception power due to distance between relay nodes and obstacles.

For achieving improvement of throughput and large circuit capacity of overall system, setting method of the relay route

becomes important. However, heretofore, there is not relay route setting method specialized radio relay transmission type cellular system in a configuration where a core node arranged a large number of micro cells, is connected to the wired backbone network and solving a problem of interference between the cells which causes problem in the cellular system.

SUMMARY OF THE INVENTION

The present invention has been worked out in view of the problem set forth above. It is therefore an object of the present invention to provide a radio network, a relay node, a core node, a relay transmission method used in the same and a program thereof, which can select a route of minimum path loss among entire relay routes and can set the relay route satisfactorily resistive against interference.

According to the first aspect of the present invention, a radio network comprises:

a core node connected to a wired network;

relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node;

a terminal station capable of transmission and reception of data packet with both of the core node and the relay node,

the relay node having total transmission loss to be minimum at least one of between relay nodes includes in a relay route

of the data packet and between the relay node and the core node is selected.

According to the second aspect of the present invention, a radio network comprises:

a core node connected to a wired network; relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node;

a terminal station capable of transmission and reception of data packet with both of the core node and the relay node,

the relay node relaying the up-link data packet to other one of up-link relay node and the core node when the up-link data packet addressed to own node is received and relaying a down-link data packet to at least one down-link relay node when the down-link data packet address to the own node is received.

According to the third aspect of the present invention, a radio network comprises:

a relay node which forgets all of update metrics corresponding to the route setting packets received in the past and relays a ner route setting packet taking the update metric corresponding to the currently received route setting packet as net metric, when a sender node identification information contained in the received route setting packet matches a current up-link receiver side relay node. According to the fourth

aspect of the present invention, a radio network comprises:

a core node connected to a wired network; relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node;

a terminal station capable of transmission and reception of data packet with both of the core node and the relay node,a radio frequency band to be used in relaying to be performed between the core node and the relay node and between the relay nodes and a radio frequency band to be used in access transmission to be performed between the core node and the terminal station and between the relay node and the terminal station, are different, and the radio frequency band to be used in relaying is higher than the radio frequency band to be used in the access transmission.

According to the fifth aspect of the present invention, a relay node relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and capable of communication with a terminal station, comprises:

an antenna for access transmission; an antenna for relaying;

a radio system for access transmission; and

a radio system for relaying,

a radio frequency band to be used in relaying to be performed between the core node and a radio frequency band to be used in access transmission to be performed between the terminal station, are different, and the radio frequency band to be used in relaying is higher than the radio frequency band to be used in the access transmission.

According to the sixth aspect of the present invention, a relay node relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and capable of communication with a terminal station,

selecting relay nodes for making a total path loss in a relay route of the data packet minimum from own node to the core node. According to the seventh aspect of the present invention, a relay node relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and capable of communication with a terminal station,

relaying the up-link data packet to other one of up-link relay node and the core node when the up-link data packet addressed to own node is received and relaying a down-link data packet to at least one down-link relay node when the down-link data packet address to the own node is received.

According to the eighth aspect of the present

invention, a relay node relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and capable of communication with a terminal station,

a radio frequency band to be used in relaying to be performed between the core node and the relay node and between the relay nodes and a radio frequency band to be used in access transmission to be performed between the core node and the terminal station and between the relay node and the terminal station, are different, and the radio frequency band to be used in relaying is higher than the radio frequency band to be used in the access transmission.

According to the ninth aspect of the present invention, a core node capable of transmission and reception of data packet with either a relay node perform radio relaying and a terminal station, and connected to a wired network, comprises: an antenna for access transmission;

an antenna for relaying;

a radio system for access transmission;

a radio system for relaying; and

a signal distributor connected to a wired backbone network,a radio frequency band to be used in relaying to be performed between the relay node and a radio frequency band to be used in access transmission to be performed between the

terminal station, are different, and the radio frequency band to be used in relaying is higher than the radio frequency band to be used in the access transmission.

According to the tenth aspect of the present invention, a core node connected to a wired network, being relayed at least one of a down-link data packet transmitted from own node and an up-link data packet directed toward own node, and capable of transmission and reception of data packet with a terminal station, transmits a route setting packet including a metric indicative of an amount providing indicia for selecting a sender identification information, an up-link receiver side relay node information and a receiver side relay node, to the relay node.

According to the eleventh aspect of the present invention, a core node connected to a wired network, being relayed at least one of a down-link data packet transmitted from own node and an up-link data packet directed toward own node, and capable of transmission and reception of data packet with a terminal station, a radio frequency band to be used in relaying to be performed between the core node and a radio frequency band to be used in access transmission to be performed between the terminal station, are different, and the radio frequency band to be used in relaying is higher than the radio frequency band to be used in the access transmission.

According to the twelfth aspect of the present invention, a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, comprises the step of: selecting the relay node having total path loss to be minimum at least one of between relay nodes includes in a relay route of the data packet and between the relay node and the core node.

According to the thirteenth aspect of the present invention, a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, comprises steps of:

relaying the up-link data packet to other one of up-link relay node and the core node when the up-link data packet addressed to own node is received and relaying a down-link data packet to at least one down-link relay node when the down-link data packet address to the own node is received.

According to the fourteenth aspect of the present invention, a relaying method comprises:

a relay node which forgets all of update metrics corresponding to the route setting packets received in the past and relays a new route setting packet taking the update metric corresponding to the currently received route setting packet as net metric, when a sender node identification information contained in the received route setting packet matches a current up-link receiver side relay node.

According to the fifteenth aspect of the present invention, a relaying method for a system including core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node,

a radio frequency band to be used in relaying to be performed between the core node and the relay node and between the relay nodes and a radio frequency band to be used in access transmission to be performed between the core node and the terminal station and between the relay node and the terminal station, are different, and the radio frequency band to be used in relaying is higher than the radio frequency band to be used in the access transmission.

According to the sixteenth aspect of the present invention, a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, comprises:

step of detecting arrival of a route setting packet including a sender node identification information, an up-link receiver side relay node information and a metric indicative of an amount providing an indicia for selecting the receiver side relay node;

step of making judgment whether the up-link receiver side relay node indicates own node or not upon detection of arrival of the route setting packet;

step of recording a node indicated by the sender node identification information contained in the route setting packet in a relay node list when judgment is made that the up-link receiver side information indicates own node;

step of taking a measured path loss upon judgment that the up-link receiver side node relay node information does not indicative own node, as path loss Ln (n is unique number of a sender node of the route setting packet) between the node

transmitting the route setting packet and the own node;

step of reading the metric Mr.n contained in the route setting packet;

step of calculating and storing an update metric from the path loss Ln and the metric Mr.n;

step of comparing the update metric Mn with the update metric corresponding to the route setting packet received in the past for making judgment whether the update metric Mn is minimum;

step of setting the update metric Mn to a metric contained in the metric of the route setting packet and registering the node indicated by the sender node identification information of the currently arrived route setting packet as the up-link receiver side relay node when the update metric Mn is judged as minimum; and

step of transmitting a route setting packet containing the transmission metric M as the metric, sender node identification information indicating identification information of own node and the up-link receiver side relay node information.

According to the seventeenth aspect of the present invention, a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the

core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, comprises:

step of detecting arrival of a route setting packet including a sender node identification information, an up-link receiver side relay node information and a metric indicative of an amount providing an indicia for selecting the receiver side relay node;

step of making judgment whether the up-link receiver side relay node indicates own node or not upon detection of arrival of the route setting packet;

step of recording a node indicated by the sender node identification information contained in the route setting packet in a relay node list when judgment is made that the up-link receiver side information indicates own node;

step of taking a measured path loss upon judgment that the up-link receiver side node relay node information does not indicative own node, as path loss Ln (n is unique number of a sender node of the route setting packet) between the node transmitting the route setting packet and the own node;

step of reading the metric Mr.n contained in the route setting packet;

step of calculating and storing an update metric from

the path loss Ln and the metric Mr.n;

step of making judgment whether the sender node identification information contained in the currently received route setting packet matches with the current up-link receiver side relay node information or not;

step of forgetting all stored update metrics when the sender node identification information contained in the currently received route setting packet matches with the current up-link receiver side relay node information;

step of comparing the update metric corresponding to the route setting packet received in the past and the currently obtained update metric Mn when the sender node identification information contained in the currently received route setting packet does not match with the current up-link receiver side relay node information;

step of setting the update metric Mn to a metric contained in the metric of the route setting packet and registering the node indicated by the sender node identification information of the currently arrived route setting packet as the up-link receiver side relay node when all of the update metrics are forgotten or when the update metric Mn is judged as minimum; and

step of transmitting a route setting packet containing the transmission metric M as the metric, sender node

identification information indicating identification information of own node and the up-link receiver side relay node information.

According to the eighteenth aspect of the present invention, a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, comprises:

step of detecting arrival of a route setting packet including a sender node identification information, an up-link receiver side relay node information and a metric indicative of an amount providing an indicia for selecting the receiver side relay node;

step of making judgment whether the up-link receiver side relay node indicates own node or not upon detection of arrival of the route setting packet;

step of recording a node indicated by the sender node identification information contained in the route setting packet in a relay node list when judgment is made that the up-link receiver side information indicates own node;

step of taking a measured path loss upon judgment that

the up-link receiver side node relay node information does not indicative own node, as path loss Ln (n is unique number of a sender node of the route setting packet) between the node transmitting the route setting packet and the own node;

step of reading the metric Mr.n contained in the route setting packet;

step of calculating and storing an update metric from the path loss Ln and the metric Mr.n;

step of comparing the update metric Mn with the update metric corresponding to the route setting packet received in the past for determining a sender node m (m is unique number of node) having minimum metric;

step of making judgment whether the sender node m is the same as the current up-link receiver side relay node and $n \neq m$;

step of setting the update metric Mn to a metric contained in the metric of the route setting packet and registering the node indicated by the sender node m as the up-link receiver side relay node when the sender node m is not the same as the current up-link receiver side relay node or $n = m$; and

step of transmitting a route setting packet containing the transmission metric M as the metric, sender node identification information indicating identification information of own node and the up-link receiver side relay

node information.

According to the nineteenth aspect of the present invention, a program of a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, the program being executed by a computer for implements the step of:

selecting the relay node having total path loss to be minimum at least one of between relay nodes includes in a relay route of the data packet and between the relay node and the core node.

According to the twentieth aspect of the present invention program of a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, the program being executed by a computer for implements the step of:

relaying the up-link data packet to other one of up-link

relay node and the core node when the up-link data packet addressed to own node is received and relaying a down-link data packet to at least one down-link relay node when the down-link data packet address to the own node is received.

According to the twenty-first aspect of the present invention, a program of a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, the program being executed by a computer for implements the step of:

step of detecting arrival of a route setting packet including a sender node identification information, an up-link receiver side relay node information and a metric indicative of an amount providing an indicia for selecting the receiver side relay node;

step of making judgment whether the up-link receiver side relay node indicates own node or not upon detection of arrival of the route setting packet;

step of recording a node indicated by the sender node identification information contained in the route setting packet in a relay node list when judgment is made that the up-link

receiver side information indicates own node;

step of taking a measured path loss upon judgment that the up-link receiver side node relay node information does not indicative own node, as path loss Ln (n is unique number of a sender node of the route setting packet) between the node transmitting the route setting packet and the own node;

step of reading the metric Mr.n contained in the route setting packet;

step of calculating and storing an update metric from the path loss Ln and the metric Mr.n;

step of comparing the update metric Mn with the update metric corresponding to the route setting packet received in the past for making judgment whether the update metric Mn is minimum;

step of setting the update metric Mn to a metric contained in the metric of the route setting packet and registering the node indicated by the sender node identification information of the currently arrived route setting packet as the up-link receiver side relay node when the update metric Mn is judged as minimum; and

step of transmitting a route setting packet containing the transmission metric M as the metric, sender node identification information indicating identification information of own node and the up-link receiver side relay

node information.

According to the twenty-second aspect of the present invention, a program of a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, the program being executed by a computer for implements the step of:

step of detecting arrival of a route setting packet including a sender node identification information, an up-link receiver side relay node information and a metric indicative of an amount providing an indicia for selecting the receiver side relay node;

step of making judgment whether the up-link receiver side relay node indicates own node or not upon detection of arrival of the route setting packet;

step of recording a node indicated by the sender node identification information contained in the route setting packet in a relay node list when judgment is made that the up-link receiver side information indicates own node;

step of taking a measured path loss upon judgment that the up-link receiver side node relay node information does not

indicative own node, as path loss Ln (n is unique number of a sender node of the route setting packet) between the node transmitting the route setting packet and the own node;

step of reading the metric Mr.n contained in the route setting packet;

step of calculating and storing an update metric from the path loss Ln and the metric Mr.n;

step of making judgment whether the sender node identification information contained in the currently received route setting packet matches with the current up-link receiver side relay node information or not;

step of forgetting all stored update metrics when the sender node identification information contained in the currently received route setting packet matches with the current up-link receiver side relay node information;

step of comparing the update metric corresponding to the route setting packet received in the past and the currently obtained update metric Mn when the sender node identification information contained in the currently received route setting packet does not match with the current up-link receiver side relay node information;

step of setting the update metric Mn to a metric contained in the metric of the route setting packet and registering the node indicated by the sender node identification information

of the currently arrived route setting packet as the up-link receiver side relay node when all of the update metrics are forgotten or when the update metric Mn is judged as minimum; and

step of transmitting a route setting packet containing the transmission metric M as the metric, sender node identification information indicating identification information of own node and the up-link receiver side relay node information.

According to the twenty-third aspect of the present invention, a program of a relaying method for a radio network including a core node connected to a wired network, relay nodes each relaying at least one of a down-link data packet transmitted from the core node and an up-link data packet directed toward the core node, and a terminal station capable of transmission and reception of data packet with both of the core node and the relay node, the program being executed by a computer for implements the step of:

step of detecting arrival of a route setting packet including a sender node identification information, an up-link receiver side relay node information and a metric indicative of an amount providing an indicia for selecting the receiver side relay node;

step of making judgment whether the up-link receiver side

relay node indicates own node or not upon detection of arrival of the route setting packet;

step of recording a node indicated by the sender node identification information contained in the route setting packet in a relay node list when judgment is made that the up-link receiver side information indicates own node;

step of taking a measured path loss upon judgment that the up-link receiver side node relay node information does not indicative own node, as path loss Ln (n is unique number of a sender node of the route setting packet) between the node transmitting the route setting packet and the own node;

step of reading the metric Mr.n contained in the route setting packet;

step of calculating and storing an update metric from the path loss Ln and the metric Mr.n;

step of comparing the update metric Mn with the update metric corresponding to the route setting packet received in the past for determining a sender node m (m is unique number of node) having minimum metric;

step of making judgment whether the sender node m is the same as the current up-link receiver side relay node and $n \neq m$;

step of setting the update metric Mn to a metric contained in the metric of the route setting packet and registering the

node indicated by the sender node m as the up-link receiver side relay node when the sender node m is not the same as the current up-link receiver side relay node or n = m; and

step of transmitting a route setting packet containing the transmission metric M as the metric, sender node identification information indicating identification information of own node and the up-link receiver side relay node information.

Namely, the radio network according to the present invention determines the core node among node groups deployed in certain area, connects the core node with the backbone network, and connects the core node with wireless. The nodes other than the core node relays up-link data to the core node or relays the down-link data transmitted from the core node.

By this, upon connecting the node groups with the backbone network, only the core node and the backbone network are connected by wire line to permit reduction of installation cost of the wire line. Also, since connection of the node groups is established by radio, service area can be easily expanded.

The core node broadcasts the relay route setting packet. The relay node measures the pathpath loss between the node transmitted the packet and the own node by reception of the relay route setting packet. At the same time, with making reference to the metric contained in the packet, the receiver

side relay station is selected so that the path loss becomes minimum by a sum of the measured path loss and the metric. Here, the metric represents a total path loss from the core node to the node transmitted the relay route setting packet.

Each base station autonomously perform the foregoing operation. Therefore, receiver side in relay to have the minimum path loss can be selected over the relay route and can establish the relay route resistive against interference which can be a significant problem in the cellular system.

Also, by using the path loss as metric, stable relay route not depending upon interference power which is variable depending upon traffic, can be certainly established. On the other hand, even when the frequency band is different, difference of the path loss is considered to be small, in general. Therefore, even when different frequency bands are used in up- and down-links, appropriate relay route can be obtained.

BRIEF DESCRIPTION OF THE DRAWINGS

The present invention will be understood more fully from the detailed description given hereinafter and from the accompanying drawings of the preferred embodiment of the present invention, which, however, should not be taken to be limitative to the invention, but are for explanation and understanding only.

In the drawings:

Fig. 1 is a diagrammatic illustration of one embodiment of a cellular system according to the present invention;

Fig. 2 is an illustration showing one example of a structure of a route setting packet;

Fig. 3 is a flowchart showing one example of a relay route setting process to be executed in each relay node in one embodiment of the present invention;

Fig. 4 is a flowchart showing one example of the relay route setting process to be executed in a core node in one embodiment of the present invention;

Fig. 5 is an illustration showing one example of a structure of an up-link data packet;

Fig. 6 is a flowchart showing one example of a transmission process of data packet in up-link;

Fig. 7 is an illustration showing a relay node list;

Fig. 8 is an illustration showing a data structure of a down-link data packet;

Fig. 9 is a flowchart showing one example of a down-link relay transmission process in one embodiment of the present invention;

Fig. 10 is a flowchart showing one example of a receiving operation of a terminal station in one embodiment of the present invention;

Fig. 11 is an illustration showing one example of a relay

route in the case where one embodiment of the relay route setting method according to the present invention is employed;

Fig. 12 is an illustration for comparing one embodiment of the relay route setting method according to the invention and a minimum hop number relay route setting method;

Fig. 13 is a flowchart showing one example of the relay route setting process to be executed in each relay node in another embodiment of the present invention;

Fig. 14 is a block diagram showing a construction of the node to be used in one and another embodiments of the present invention;

Fig. 15 is a block diagram showing a construction of the relay node to be used in one and another embodiments of the present invention;

Fig. 16 is a block diagram showing a construction of the core node to be used in one and another embodiments of the present invention;

Fig. 17 is a flowchart showing an example of another process in a part of the relay route setting process in one and another embodiments of the present invention;

Fig. 18 is a flowchart showing one example of the relay route setting process to be executed in each relay node in a further embodiment of the present invention;

Fig. 19 is a flowchart showing an example of another process

in a part of the relay route setting process in another embodiment of the present invention; and

Fig. 20 is an illustration diagrammatically showing the conventional cellular system.

DESCRIPTION OF THE PREFERRED EMBODIMENT

The present invention will be discussed hereinafter in detail in terms of the preferred embodiment of the present invention with reference to the accompanying drawings. In the following description, numerous specific details are set forth in order to provide a thorough understanding of the present invention. It will be obvious, however, to those skilled in the art that the present invention may be practiced without these specific details. In other instance, well-known structure are not shown in detail in order to avoid unnecessary obscurity of the present invention.

Fig. 1 is a diagrammatic illustration showing one embodiment of a cellular system according to the present invention. In Fig. 1, the reference numeral 107 denotes a terminal station, 108 denotes a cell. A core node 103 and a wired backbone 101 are connected by wire circuit 102. Relay nodes 104 to 106 are connected to the core node 103 by radio relay.

Each of the relay nodes and the core node may be provided with directional antenna. In such case, the directional antenna

may be fixedly set directional, or, in the alternative, may be set directionality adaptively. By installing directional antenna, interference given to neighboring nodes and the terminal stations can be suppressed, which realize high system capacity in overall system.

Other relay nodes shown in Fig. 1 are also connected to the core node 103 through radio relay circuit, similarly to the relay nodes 104 to 106. Setting of relay route of the radio relay circuit is initiated by a route setting packet broadcasted from the core node 103. Namely, the relay node receiving the route setting packet broadcasted by the core node 103 newly broadcast a route setting packet to other nodes. In response to those broadcasted route setting packet sent by the relay node, other relay nodes also broadcast the route setting packet. This operation is repeated. Detail of transmission of the route setting packet will be discussed later.

Fig. 2 is an illustration showing one example of a structure of the route setting packet. In Fig. 2, the route setting packet is consisted of fields for transmitting respective of a sender node ID (identification information) A02, an up-link receiver side relay node ID A03, a metric A04 and others A01. It should be noted that order of arrangement of respective elements can be different from an example shown in Fig. 2.

The sender node ID A02 indicates an ID number of the node

transmitted the route setting packet. The up-link receiver side relay node ID A03 indicates ID number of an up-link receiver side relay node set by the node broadcasting the route setting packet. In other A01, control signal, such as pilot signal or the like, data signal such as system information and so forth, necessary for demodulation of the packet, are included. The metric A04 indicates the amount of information for each node to select an up-link receiver side relay node.

Fig. 3 is a flowchart showing one example of a relay route setting process executed in each relay node in one embodiment of the present invention, and Fig. 4 is a flowchart showing one example of the relay route setting process to be executed in a core node in one embodiment of the present invention. Updating method of the metric A04 and up-link receiver side node selecting procedure and route setting process in the code node 103 depending upon the metric amount will be discussed with reference to Figs. 1 to 4.

At first, transmission of the route setting packet is initially performed by the core node 103. The relay route setting packet broadcasted from the core node 103 is received by unspecified relay nodes 104 to 106. Namely, transmission of the route setting packet is performed by broadcasting. At this time, since no up-link receiver side relay node of the core node 103, the content of the up-link receiver side node

ID A03 may be any ID as arbitrarily set.

The metric contained in the route setting packet broadcasted from the code node 103 is set to zero. A broadcasting interval of the route setting packet may be set for regularly broadcasting, or in the alternative for broadcasting at random interval or in response to a command from a server (not shown) on the wired backbone network 101.

The relay nodes 104 to 106 check whether the route setting packet arrived or not (step S1 of Fig. 3). If the route setting packet has not arrived, the process is returned to step S1. For detection of arrival of the route setting packet in the relay nodes 104 to 106, carrier sense or the like is used. When arrival of the route setting packet is detected, the relay nodes 104 to 106 makes reference to the up-link receiver side relay node ID contained in the route setting packet for checking whether the up-link receiver side relay node ID matches with the own node ID or not (step S8 of Fig. 3).

When the receiver side relay node ID matches with the own node ID, the relay nodes 104 to 106 record the ID of the node broadcasted the route setting packet, namely the sender node ID contained in the route setting packet in the relay node list (step S9 of Fig. 3).

The relay node list is a table indicating the down-link receiver side node ID number and is constructed as shown in

Fig. 7. The relay node list is used as receiver side node list upon the down-link data packet relaying which will be discussed later. Each receiver side relay node ID contained in the relay node list can be forgotten (erased) after elapse of a given period. For instance, when new relay node is added in the cell, when existing node moved or when new building is constructed in the cell, re-establishment of the relay route becomes necessary. In order to adapt to re-established relay route, each receiver side relay node ID contained in the relay node list may be intentionally forgotten after expiration of the given period.

If judgment is made that the relay node ID does not match with the own node ID, the relay nodes 104 to 106 takes a path loss as measured at the timing of judgment, as a path loss Ln (n is specific number of the sender node of the route setting packet) between the node that broadcasted the route setting packet and the own node (step S2 of Fig. 3). Measurement of the path loss is generally performed upon reception of the packet irrespective of the content. For measurement of the path loss, reception power of the packet or the like may be used. In order to facilitate measurement of the path loss, transmission power of the route setting packet can be fixed. It should be noted that n represents a node number and the node number n is designated by the sender node ID contained in the route setting packet.

The relay nodes 104 to 106 read the metric Mr, n contained in the received route setting packet (step S3 of Fig. 3). Here, the metric Mr, n represents total of path loss in the set route. Here, the set route means a route from the sender node of the received route setting packet to the core node.

The relay nodes 104 to 106 set an update metric Mn from the transmission loss Ln measured at step S2 and the metric Mr.n. Here, the update metric Mn is given as a sum of the path loss Ln and the metric Mr.n. The relay nodes 104 to 106 store the update metric Mn calculated through the foregoing process (step S4 of Fig. 3). It should be noted that, among stored update metric, the update metric maintained for a period in excess of a given period can be forgotten (erased). For instance, when new relay node is added in the cell, when new building is constructed in the cell, re-establishment of the relay route becomes necessary. In order to adapted to this, stored update metrics may be intentionally forgotten after expiration of the given period.

On the other hand, stored metric is made most recent constantly. Namely, when the update metric for the node n as sender of the route setting packet was stored in the past, the past metric is re-written by the new update metric derived at step S4.

The relay nodes 104 to 106 compares the update metric

Mn and the update metric corresponding to the route setting packet received in the past. If the newly obtained update metric Mn is not minimum (step S5 of Fig. 3), the process is returned to step S1 and route setting packet is not transmitted.

When the currently obtained update metric Mn is minimum (step S5 of Fig. 3), the relay nodes 104 to 106 set the metric to be contained in the metric A04 to the update metric Mn, and the node indicated by the sender node ID of the currently arrived route setting packet is registered as the up-link receiver side relay node (step S6 of Fig. 3). Thus, the up-link receiver side relay node is only one in each node.

The relay nodes 104 to 106 sets the metric M set as set forth above as metric and transmit the route setting packet with including the necessary information in respective of other items shown in Fig. 2 (step S7 of Fig. 3).

Upon reception of the route setting packet, the relay nodes 104 to 106 may return an acknowledgement signal for accuracy. Since the route setting packet is the control packet directed to unspecified node, the relay nodes 104 to 106 may receive the reception response signals from a plurality of nodes after transmission of the route setting packet. When the relay nodes 104 to 106 does not receive returned reception response signal at all, re-transmission of the route setting packet is performed.

On the other hand, the relay route setting process to

be executed in a core node is different from that disclosed for relay nodes. Initially, the core node 103 checks whether the route setting packet arrived or not (step S11 of Fig. 4). If the route setting packet has not arrived, the process is returned top step S11. Even in the core node 103, carrier sense or the like is used for detection of arrival of the route setting packet. When the core node 103 detects arrival of the route setting packet, the up-link receiver side relay node ID contained in the route setting packet is made reference to check whether the up-link receiver side node ID matches with the own node ID (step S12 of Fig. 4).

When the up-link receiver side node ID matches with the own node ID, the core node 103 records the node indicated in the ID of the route setting packet transmitted by the other node, namely the sender node ID contained in the route setting packet, in the relay node list (step S13 of Fig. 4). The relay node list has the same function as those contained in the relay nodes 104 to 106. Namely, the relay node list is a table indicating number ID of the down-link receiver side relay node and respective receiver side relay node contained in the relay node list can be forgotten (erased) after expiration of a given period.

Next, relay transmission of the data packet will be discussed. Fig. 5 is an illustration showing one example of

the structure of the up-link data packet. In Fig. 5, the up-link receiver side data packet is consisted of fields respectively transmitting a receiver side relay node ID B02, a sender side relay node ID B03, a source terminal ID B04, a data B05 and other B01.

In the sender side relay node ID B03, the ID of the relay node that transmitted the up-link data packet or is going to transmit the up-link data packet is set. When the terminal transmits new up-link data packet, special information indicative of a condition of the new up-link data packet other than node ID is contained in the sender side relay node ID, is transmitted.

In the field of other B01, control information, such as pilot signal for decoding, identification signal identifying up and down-links, ID number of the data packet and so forth is contained. It should be noted that the order of the structural elements shown in Fig. 5 should not be necessarily specific but can be different in any order.

Fig. 6 is a flowchart showing one example of a transmission process of data packet in up-link. Relay transmission method of the data packet in up- and down-links in one embodiment of the present invention will be discussed with reference to Figs. 5 and 6. At first, in one embodiment of the present invention, one example of the relay transmission method of the up-link

to be executed in each base station in one embodiment of the present invention will be discussed.

The up-link data packet is transmitted to the code node 103 via the relay nodes 104 to 106. The relay nodes 104 to 106 detect arrival of the up-link data packet (step S21 of Fig. 6). Here, for detection of the data packet, carrier sense or the like is used. Judgment whether the data packet of interest is up-link one or down-link is performed depending upon the control information contained in the up-link data packet shown in Fig. 5.

When a plurality of nodes relay the data packet from the same terminal, it is possible that data packets of the same content are received from a plurality of sender nodes by one node. In such case, demoudulating may be performed by selecting only up-link data packet having the highest reception quality or by combining the received signals by diversity reception technique.

If arrival of the up-link data packet is not detected, the relay nodes 104 to 106 execute setp S11 again. On the other hand, when arrival of the up-link data packet is detected, the relay nodes 104 to 106 checks whether the arrived up-link data packet is the data pack on relay or the data packet newly transmitted from a terminal 107 (step S22 of Fig. 6).

Here, upon making judgment wether the packet is the data

packet on relay, the node 104 to 106 check the sender side node ID B03 contained in the up-link data packet. For instance, if the sender side relay node ID B03 indicates the ID of the own node, it is judged that the packet is a new up-link data packet.

Here, upon making judgment wether the packet if the data packet on relay, the relay node 104 to 106 check the sender side node ID B03 contained in the up-link data packet. For instance, if the sender side relay node ID B03 is the ID of the own node, it is judged that the packet is a new up-link data packet.

Here, upon making judgment of the data packet on relay (other than new up-link data packet), the relay node 104 to 106 check the receiver side relay node ID B02 contained in the data packet. If the receiver side relay node ID B02 checked is not the ID of the own node (step S23 of Fig. 6), the process returns to step S21.

If the receiver side relay node ID B02 checked is the own node ID (step S23 of Fig. 6), the relay node 104 to 106 record the sender side relay node ID B03 in the relay node list (step S24 of Fig. 6). One example of the relay node list is shown in Fig. 7.

The relay node list is used as the receiver side node

list upon down-link data packet relaying as will be discussed later. Each receiver side relay node ID B02 contained in the relay node list can be forgotten after expiration of the given period. In the node not receiving the up-link data packet from the relay node, the relay node list becomes empty.

After recording the sender side relay node ID B03, the relay nodes 104 to 106 relay the data packet to the up-link receiver side relay node set in the relay route setting process set forth above (step S25 of Fig. 6). After relaying, the process returns to step S21.

Upon transmission of the up-link data packet, transmission power of the data packet can be controlled so that the data packet may be received in the relay node or the receiver side relay node with a given reception power or a given reception quality.

On the other hand, when judgment is made that the arrived up-link data packet is not on relaying but is newly generated from the terminal station 107 (step S22 of Fig. 6), the relay nodes 104 to 106 relay the data packet toward the receiver side relay node (step S25 of Fig. 6). It should be noted that recoding operation to the relay node list upon the up-link data packet shown in Fig. 6 is also performed upon relay route setting packet transmission set forth above. For the reason of reduction of the process load, recording operation may not be implemented

upon up-link data packet transmission.

The relaying of the up-link data packet in the core node is substantially the same as that in the relay node shown in Fig. 6 except for step S25. In the core node, instead of relaying the up-link data packet to the receiver side relay node, the up-link data packet is transmitted to the wired backbone network.

Fig. 8 is an illustration showing a data structure of the down-link data packet. In Fig. 8, the data packet is consisted of fields respectively transmitting a receiver side relay node IS C02, a sender side relay node ID C03, a destination terminal ID C04, a data C05 and other C01.

In the field for the sender side relay node ID C03, ID of the core node 103 transmitting the down-link data packet or the relay nodes 104 to 106 is set. When a plurality of receiver side relay nodes are present, a plurality of receiver side relay node ID C02 are also provided. On the other hand, the receiver side node ID C02 can set only the individual node ID but also dedicated ID indicative of all nodes included in the relay node list. In the field for other C01, control information, such as pilot signal for demoudulating, identification signal of up- and down-link and packet ID number and so forth can be contained. It should be noted that the order of the structural element shown in Fig. 8 should not be specific but can be different.

Fig. 9 is a flowchart showing one example of a down-link relay transmission process in one embodiment of the present invention. Discussion will be given for one example of the down-link data packet relaying process in one embodiment of the present invention will be discussed with reference to Figs. 8 and 9. It should be noted that the process shown in Fig. 9 is implemented in each relay nodes 104 to 106.

The relay nodes 104 to 106 monitors arrival of the down-link data packet. If no down-link data packet newly arrives (step S31 of Fig. 9), the process is returned to step S31. Detection of arrival of the down-link data packet is implemented by carrier sense or the like.

When the down-link data packet newly arrives (step S31 of Fig. 9), the receiver side relay node contained in the down-link data packet is read. If the receiver side node ID does not match with the own node ID (step S32 of Fig. 9), the process returns to step S31 not to perform relaying of the reception data packet.

If the receiver side relay node ID matches with the own node ID (step S32 of Fig. 9), the relay nodes 104 to 106 make reference to the relay node list produced upon relaying of the up-link data packet or relaying of the route setting packet to select part of all of nodes contained in the relay node list to set as the receiver side relay node for the data packet (step

S33 of Fig. 9).

Upon setting all nodes, particular identification number dedicated therefor is set as the sender ID C04. The relay nodes 104 to 106 relays the data packet after setting of the receiver side relay node (step S34 of Fig. 9).

Upon transmission of the down-link data packet, transmission power of the data packet can be controlled so that the data packet may be received in the relay node or the receiver side relay node with a given reception power or a given reception quality.

Relaying of the down-link data packet in the core node can be the same as the method in the relay node shown in Fig. 9.

Fig. 10 is a flowchart showing one example of a receiving operation of a terminal station 107 in one embodiment of the present invention. One example of receiving operation of the terminal station 10 in one embodiment of the present invention will be discussed with reference to Fig. 10.

The terminal station 107 detects arrival of the down-link data packet by carrier sense or the like to return to step S41 if arrival of the data packet is not detected (step S41 of Fig. 10). When arrival of the data packet is detected (step S41 of Fig. 10), the destination terminal ID contained in the down -link data packet shown in Fig. 8 is read out. If the desitination

terminal ID does not match with the own terminal ID (step S42 of Fig. 10), the process returns to step S41.

If the destination terminal ID matches with the own terminal ID (step S42 of Fig. 10), the terminal station 107 performs reception process of the data contained in the data packet (step S43 of Fig. 10). Then, process returns to step S41.

Figs. 11 and 12 are illustrations showing one example of the relay route set by relay route setting in one embodiment of the present invention. Fig. 11 is an illustration showing one example of a relay route in the case where one embodiment of the relay route setting method according to the present invention is employed. In Fig. 11, a large dot 201 represents the core node, and small dots 202, 204, 205 and so on represent relay nodes, and 203 denotes relay routes. At the relay node 202, no down-link receiver side relay node is present, which means that the relay node list becomes empty in the relay node.

Fig. 12 is an illustration for comparing one embodiment of the relay route setting method according to the invention and a minimum hop routing method. In Fig. 12, 301 denotes a wired backbone network, 302, 303 and 304 are cells covered by the core nodes, 309 and elliptic areas not identified by reference numeral are nodes covered by nodes other than core node. The reference numerals 310, 311, 312 are wire lines connecting the

core nodes and the wired backbone network 301.

One example of the relay route obtained by the relay route setting method by one embodiment of the present invention is shown as radio relay circuits 307, 305 and 306. On the other hand, for comparison, minimum hop routihg, number relaying, with minimum number of relay nodes, provides the radio relay route shown by reference numeral 308 for example.

By using the relay route setting method of one embodiment of the present invention, selecting the route to have the minimum path loss in the overall relay route can be realized. Then, relay route strong against interference can be realize.

On the other hand, in case of the minimum hop route shown in Fig. 12, number of relay stations becomes smaller in comparison with the route setting method in one embodiment of the present invention. However, considering overall relay route, total path loss becomes greater than that in one embodiment of the present invention. Therefore, reliability of the overall radio relay route becomes low. In the relaying method in one embodiment of the present invention, the radio relaying route having high reliability can be certainly established, which renders higher throughput in comparison with minimum hop route method.

In the down-link, using a relationship between the sender side relay node and the receiver side relay node formed upon

up-link packet relay, the node ~~served as sender~~ side relay node in up-link is selected as the receiver side relay node.

By using the path loss as the metric, stable relay route independent of interference power which is variable depending upon amount of traffic, can be certainly established. On the other hand, since difference of path loss for different frequency bans is generally considered to be small appropriate relay route can be established if different frequency bands are used in the up- and down-link.

Only core node is connected to the backbone network by wire, and connection between other node group and the backbone network is automatically established by radio, therefor installation cost of the wire line can be reduced. Also, since the node groups are connected by radio, service area can be easily expanded. Furthermore, each node is not constrained by wire, re-arrangement of node location can be done easily.

When the terminal moves between the nodes belonging the same core node, it becomes unnecessary to access the mobile control station or the like in the wired backbone network, and thus high speed hand-over becomes possible.

As can be clear from the example shown in Fig. 11, when the up-link data packet transmitted by the terminal is received by a plurality of nodes, the same up-link data packet is relayed through a plurality of relay routes. As can be appreciated from

the example shown in Fig. 11, in relaying of the up-link data packet, the relay routes are inherently coupled together on the way of relaying at certain node, diversity effect can be obtained by selecting one having good quality upon receiving the data packet or by combining the same up-link data packet in the node where the relay routes are coupled.

In the cellular system premised in one embodiment of the present invention, a terminated node of the up-link and a start node of the down-link of the relay node are both core node. Therefore, in comparison with route setting in an adhoc network or the like, memory amount and complexity of the route setting .

Fig. 13 is a flowchart showing one example of the relay route setting process to be executed in each relay node in another embodiment of the present invention. Another embodiment of the present invention has similar construction as the one embodiment of the cellular system of the present invention shown in Fig. 1, and the structure of the route setting packet to be used in the operation thereof is also similar to the structure of one embodiment of the route setting packet according to the present invention shown in Fig. 2. With reference to Figs. 1, 2 and 13, updating method of the metric A01 and receiver side relay node selection procedure based on the metric in another embodiment of the present invention will be discussed.

Transmission of the route setting packet is initially

performed by the core node 103. The relay route setting packet transmitted from the core node 103 to the unspecified relay nodes 104 to 106.

The metric contained in the route setting packet transmitted by the core node 103 is set to zero. Transmission transmission interval of the route setting packet is set to be constant, to be random or is set in response to the command from the server on the wired backbone network.

At first, the relay nodes 104 to 106 checks whether the route setting packet arrived or not (step S51 of Fig. 13). If the route setting packet does not arrive, the process returns to step S51.

The relay nodes 104 to 106 uses carrier sense or the like for detection of arrival of the route setting packet. When arrival of the route setting packet is detected (step S51 of fig. 13), the relay node 104 to 106 makes reference to the up-link receiver side relay node ID included in the route setting packet to make judgment whether the up-link receiver side relay node ID matches with the own node ID or not (step S60 of Fig. 13).

When the receiver side relay node ID matches with the own node ID, the relay nodes 104 to 106 record the ID within the route setting packet which the node broadcasted, namely the sender node ID contained in the route setting packet in the relay node list (step S61 of Fig. 3).

The relay node list is a table indicating the down-link receiver side node index and is constructed as shown in Fig. 7. The relay node list is used as receiver side node list for the down-link data packet relaying which was discussed befor, and will be discussed later too. Each receiver side relay node ID contained in the relay node list can be forgotten (erased) after elapse of a given period. For instance, when new relay node is added in the cell, when existing node moved or when new building is constructed in the cell, re-establishment of the relay route becomes necessary. In order to adapt to re-established relay route, each receiver side relay node ID contained in the relay node list may be intentionally forgotten after expiration of the given period.

If judgment is made that the relay node ID does not match with the own node ID, the relay nodes 104 to 106 takes a path loss as measured at the timing of judgment, as a path loss Ln (n is specific number of the sender node of the route setting packet) between the node that broadcasted the route setting packet and the own node (step S62 of Fig. 3). Measurement of the path loss is generally performed upon reception of the packet irrespective of the content. For measurement of the path loss, reception power of the packet or the like may be used. In order to facilitate measurement of the path loss, transmission power of the route setting packet can be fixed. It should be noted

that n represents a node number and the node number n is designated by the sender node ID contained in the route setting packet.

The relay nodes 104 to 106 read the metric Mr, n contained in the received route setting packet (step S53 of Fig. 3). Here, the metric Mr, n represents total of path loss in the set route. Here, the set route means a route from the sender node of the received route setting packet to the core node.

The relay nodes 104 to 106 set an update metric Mn from the transmission loss Ln measured at step S52 and the metric Mr.n. Here, the update metric Mn is given as a sum of the path loss Ln and the metric Mr.n. The relay nodes 104 to 106 store the update metric Mn calculated through the foregoing process (step S54 of Fig. 3).

After setting the update metric Mn, if the sender node ID contained in the currently received route setting packet matches with the current up-link receiver side node ID of own node (step S55 of Fig. 13), the metric M to be contained in the field of the metric A04 with the update metric Mn and the node indicated by the sender node ID of the currently arrived route setting packet is registered as the receiver side relay node (step S58 of Fig. 13) after forgetting all of the stored update metric (step S56 of Fig. 13), Thus, up-link receiver side node is only at each node.

The relay nodes 104 to 106 sets the metric M set as set

forth above as metric and transmit the route setting packet with including the necessary information in respective of other items shown in Fig. 2 (step S59 of Fig. 13).

On the other hand, when the sender node ID contained in the currently received route setting packet does not match with the current up-link receiver side relay node ID (step S55 of Fig. 13), the update metric corresponding to the route setting packet received in the past is compared with the newly obtained update metric Mn (step S57 of Fig. 13).

If the update metric Mn is minimum, the metric M to be contained in the field of the metric A04 is set with the update metric Mn, and the node indicated by the sender node ID of the currently arrived route setting packet is registered as the up-link receiver side relay node (step S58 of Fig. 13). Thus, up-link receiver side node is only at each node.

The relay nodes 104 to 106 set the transmission metric M set as set forth above as metric to transmit the route setting packet with containing necessary information in each item shown in Fig. 2 (step S59 of Fig. 13). It should be noted that if the update metric is minimum, the process returns to step S51.

By forgetting update metric stored in each node and the receiver side relay node ID contained in the relay node list, re-establishment of the relay route becomes possible in case that modification of the relay route becomes necessary due to

variation of path loss between the nodes and addition or deletion or the like of the relay node.

On the other hand, in the relay node, concerning the case where the route setting packet transmitted from the current up-link receiver side relay node is received, all of stored update metrics corresponding to the past received route setting packets are forgotten. And transferring new update metric calculated from the metric contained in the route setting packet to the route setting packet as new metric is performed to promote updating of the route setting. Thus, it becomes possible to adapt for variation of the path loss in the current relay route.

In the cellular system premised in another embodiment of the present invention, since the relay node is stationary arranged as infrastructure, in comparison with the adhoc network, in which the moving terminal also serves as relay station, more stable communication can be performed. On the other hand, in the present invention, since both of the terminated node of the up-link and the start node of the down-link on the relay route are core node. Therefore, in comparison with route setting in an adhoc network or the like, memory amount and complexity of the route setting method can be reduced.

By controlling the transmission power of the uplink data packet or the downlink data packet, the interfarence to the nodes and the terminal station can be reduced. As a result,

the capacity of the whole system can be improved.

In the radio network in the present invention, the radio frequency band to be used for relaying between the core node and the relay node and between the relay nodes and the radio frequency band to be used for access transmission to be performed between the core node and the terminal station or between the relay node and the terminal station, can be the same or different. In transmission between semi-stationary arranged core node and relay node relative high frequency can be used, such as submillimeter wave, millimeter wave or the like. Because there are relatively large margin in frequency resource and it becomes possible to use radio wave having straight transmission characteristrics. On the other hand, in transmission between the moving terminal station and the core node or between the relay nodes, relatively low frequency band, such as microwave band, is used to permit large capacity relaying and to permit access transmission capable of non line-of-sight communication.

In the embodiment set forth above, it is premised that radio wave is used in relaying and access transmission, it is possible to use infrared light, light or the like in place of the radio wave.

Fig. 14 is a block diagram showing a construction of the node to be used in one and another embodiments of the present invention. In Fig. 14, the shown nodes are provided with

directional antennas 11 to 1n. The directional antennas 11 to 1n are connected to an antenna controller 1 through signal lines 21 to 2n to permit control of respective directionality by the antenna controller 1. Through the signal lines 21 to 2n, transmission of the transmission and reception signals and transmission of the control signal instructing the antenna direction are performed.

The antenna controller 1 is connected to a transceiver 2 via a signal line 30 for transmission of the data signal and the control signal therethrough. The antenna controller 1 performs selecting control or combining control of the transmission and reception antenna. IN the transceiver, demodulation of the received data signal, modulation of the transmission signal and so forth are unitarily processed.

In the construction shown in Fig. 14, a plurality of antennas are selected and used for one transceiver 2. However, it is also possible to perform plurality of transmission simultaneously by using the independent transceiver for each antenna.

As set forth above, by using the directional antennas 11 to 1n, it becomes possible to compensate significant distance attenuation which can be caused in using high frequency, such as millimeter wave and to obtain large gain.

Since the present invention is directed to a network

adaptively setting routes on the basis of the metrics from peripheral nodes, large gain can be obtained with respect to the node selected as route and can reduce interference for the nodes located nearby but out of the route by directing the transmission direction of the directional antennas 11 to 1n toward the node selected as the route. On the other hand, concerning the route setting packet, a non-directional antenna may be used for widely broadcasting to adjacent nodes.

Fig. 15 is a block diagram showing a construction of the relay node to be used in one and another embodiments of the present invention. In Fig. 15, a construction of the relay node 3 is shown in the case where different radio frequency bands are used for access transmission and relaying. In this case, the relay node 3 is provided with antenna for access 32 and a antenna for relay 31, and radio system for access 33 and radio for relay system 34.

The radio system for access 33 and the radio system for relay 34 includes modulator and demodulator, encoding and decoding device and so forth. The radio system for access 33 and the radio system for relay 34 can exchange signal therebetween. The radio system for relay 34 performs relaying on the basis of the route set by any one of the one embodiment of the invention and another embodiment of the invention. It should be noted that, in the route setting portion 36, setting of the route

by any one of the one embodiment of the invention and another embodiment of the invention is performed by a program stored in a storage medium 37.

At first, discussion will be given for a manner of transmission of an up-link traffic directed from the terminal station to the wired backbone network. The up-link traffic from the terminal station generated in the cell of the own node is received by the access transmission antenna 32, is processed by the accessing radio system 33 and is input to the radio system for relay 34. The radio system for relay 34 transmits the up-link traffic toward the up-link receiver side relay node using the antenna controller 35 and the antenna for relay 31.

Next, discussion will be given for operation for relaying the packet from the terminal station generated in the cell of other relay node. Namely, in the following discussion, operation in the relay node included in the route to the core node in the relay node, in which the terminal station generated the packet belongs is indicated. When the packet from the terminal station is to be relayed, at first, the up-link data packet is received by the antenna for relay 31. The received signal is input to the radio system for relay 34. Then, operation shown in Fig. 6 is performed. When relaying of the packet is decided, the radio system for relay 34 transmits the up-link traffic toward the up-link receiver side relay node using the

antenna controller 35 and the antenna for relay 31.

Subsequently, discussion will be given for a manner of transmission of the down-link traffic from the wired backbone network to the terminal station. The down-link traffic is at first received by the antenna for relay 31 and is input to the radio system for relay 34 through the antenna controller 35. The radio system for relay 34 inputs the down-link traffic to the radio system for acess 33 when the received down-link traffic is address to the terminal in the cell of own node, and otherwise, the down-link receiver side relay node is determined on the basis of the destination of the downlink packet and transmits the packet through the antenna controller 35 and the antenna for relay 31. The radio system for acess 33 is responsive to input of the down-link traffic to transmit the same to the terminal station through the antenna for acess 32.

Fig. 16 is a block diagram showing a construction of the core node to be used in one and another embodiments of the present invention. In Fig. 16, there is shown a construction of the core node in the case where different radio frequency bands for access transmission and relaying. The core node 4 has similar construction as the relay node 3 but is differentiated from the relay node as connected to the wired backbone network. A signal distributor 46 in the core node 4 is connected to the wired backbone network 40, the radio system for acess 43 and

the radio system for relay 44, respectively.

The radio system for relay 44 performing relaying based on the route set by any one of the one embodiment of the invention and another embodiment of the invention by a route setting portion 48. It should be noted that, in the route setting portion 46, setting of the route by any one of the one embodiment of the invention and another embodiment of the invention is performed by a program stored in a storage medium 49.

At first, discussion will be given for a manner of transmission of an up-link traffic directed from the terminal station to the wired backbone network 40. The up-link traffic from the terminal station generated in the cell is received by the access antenna for access 42, is processed by the radio system for access 43 and is input to the signal distributor 46. The signal distributor 46 transmits the up-link traffic to the wired backbone network 40.

Next, discussion will be given for operation for relaying the packet from the terminal station generated in the cell of other relay node. Namely, in the following discussion, operation in the case where the packet from the relay node, in which the terminal transmitted the packet belongs, reaches

core node is indicated. When the packet from the terminal station is to be relayed, at first, the up-link data packet

is received by the antenna for relay 41. The received signal is input to the relaying radio system 44. When confirmation is made that the received packet is up-link packet, the packet is transmitted to the wired backbone network through the signal distributor 46.

Subsequently, discussion will be given for a manner of transmission of the down-link traffic from the wired backbone network 40 to the terminal station. The down-link traffic is at first input to the signal distributor 46. The signal distributor 46 makes judgment of the input down-link traffic. If the down-link traffic is addressed to the terminal station within the cell of the own node, the traffic is input to the radio system for access 43, and otherwise, is input to the radio system for relay 44. When the down-link traffic is input to the radio system for access 43, the traffic is transmitted to the terminal station through the antenna for access 42. Then the down-link traffic is input to the radio system for relay 44, the traffic is transmitted to the other relay node throuth the antenna for relay 41.

In one embodiment of the present invention, upon updating of path loss as metric, new metric is constantly derived by adding the measured transmission loss Ln between the node that transmitted the route setting packet and the own node to the metric Mr,n contained in the route setting packet. However,

it is also possible to generate the update metric by multiplying the received metric and the measured path loss by a weighting coefficient having a value in a range of 0 to 1. Namely, as weighting coefficient α, new update metric is derived by (Mr,n) x α + Ln x (1 - α). With taking the value of α as 0.5, it becomes equivalent to the case where weighting is not provided and is effective for reducing transmission power in the overall system. On the other hand, by setting the value of α zero, instead of path loss from the core node, only path loss between the immediately adjacent node is considered to be effective for lowering of the transmission power in each node. Thus, by performing weighting upon updating of the metric, characteristics of the route can be varied flexibly.

On the other hand, while one embodiment of the present invention has been discussed using only path loss as metric, it is possible to use two kinds of metrics. Namely, by providing two metrics of the first metric and the second metric, when the first metrics for multiple routes are the same, judgment is made by comparing the second metric for respective routes. For example, the first metric is set a total of hop count and the second metric is set a total of the path loss. Then, hop count as the first metric is the same and minimum, one route having smaller path loss as second metric is taken as up-link side route to transmit new route setting packet and to set the

up-link receiver side relay node.

By this, more detailed route setting becomes possible to permit selection of the route having small path loss with restricting increase of delay to be caused by increasing of the hop count to restrict interference.

As set forth above, by using two kinds of metrics, characteristics of the generated route can be defined in detail to make the network close to those expected by a designer.

On the other hand, when two kinds of metrics are used, upon making judgment as the same metric, width is provided for the judgment reference to make the metric falling within a defined range as the same metric. Namely, metrics to be considered as comparable extent are considered as the same to relay on judgment of the second metric. Discussion will be given hereinafter in the case where the path loss is classified into reference per 10 dB (0 to 10 dB, 10 to 20 dB, ...) with taking total of path loss as first metric and total hop count as second metric, for example.

Assuming that three routes, route A, route B and route C respectively having first and second metrics, route A = (81 dB, 3 hops), route B = (85dB, 2 hops) and route C = (103 dB, 2 hops), are present. At this time, at first, the first metrics are compared. In the shown case, while magnitudes of the path losses <u>per se</u> are different between the route A and the route

B, in view of the reference of 10 dB unit, both falls within a reference value range of 80 to 90 dB. Therefore, the first metrics in the route A and the route B can be regarded as same. The route C has greater magnitude of path loss than those of the routes A and B even in view of the reference range of 10 dB unit. Therefore, the route C is not selected. Concerning the routes A and B judged to have the same metric, two metrics are compared. Since the route B has smaller hop count, i.e. second metric, the route B is selected as the route.

As set forth, when two metrics are used, by providing width in the reference range upon making judgment of large and small of the metric, two metrics are appropriately used to permit generation of appropriate route.

By combining a method for multiplying weighting coefficient to the metric and a method using two kinds of metrics, the following operation can be considered for example. Among two kinds of metrics, the hop count is taken as first metric. On the other hand, as second metric, with taking the path loss set forth above is taken as reference. Upon updating of metric, with taking weighting coefficient α for hop count as first metric to be $\alpha = 0.5$, and 0.5 is constantly added upon updating of metric (since hop count is constantly increased by one, 1 x 0.5 = 0.5 is used), and as weighting coefficient β for path loss as second metric to be $\beta = 0$, the measured path loss <u>per</u>

se is used as the update metric.

Namely, in Figs. 3 and 13 showing the relay route setting process, the metric Mr,n to be read at step S3 or step S53 represents a total of the hop number from the core node as the first metric and the path loss as second metric. Updating of metric at step S4 or S54 is performed respectively for the first metric and the second metric to obtain the update metric Mn.

Comparing the first update metric obtained from this result and the first update metric corresponding to the route setting packet from other node received in the past, the metric M is newly set in the update metric Mn when the newly obtained first update metric is minimum or the same as the minimum value of the first update metric corresponding to the route setting packet from other node received in the past, and when the second update metric is smaller than the second metric corresponding to the route having the minimum value of the first update metric, and the node indicated by the sender node ID of the currently received set packet is registered as up-link receiver side relay node.

Namely, in Figs. 3 and 13, upon judgment of minimum metric at step S5 or step S7, two kinds of metrics are used as set forth above. Fig. 17 shows a flowchart showing the operation. Steps S71 to S73 shown in Fig. 17 are replacement of step S5 or S57 in Fig. 3 or Fig. 13. For example, it is assumed as

combination of the first metric and the second metric, three routes, route A = (3 hops, 100 dB), route B (3 hops, 91 dB) and route C (4 hops, 85 dB) are present. At first, the route C having large hop count as the first metric is eliminated. Subsequently, by comparing the second metrics, the route B is selected.

Fig. 18 is a flowchart showing one example of the relay route setting process to be executed in each relay node in a further embodiment of the present invention. The further embodiment of the present invention has similar construction as one embodiment of the cellular system shown in Fig. 1 according to the present invention. A structure of the route setting packet used in the operation of the further embodiment of the present invention is similar to the structure of the route setting packet in one embodiment of the present invention shown in Fig. 2. Updating method of the metric A04 and receiver side relay node selection procedure in the amount of the metric in the further embodiment of the present invention will be discussed with reference to Figs. 1, 2 and 18.

Transmission of the route setting packet is initially performed by the core node 103. The relay route setting packet transmitted from the core node 103 is relayed to unspecified relay nodes 104 to 106.

The metric contained in the route setting packet

transmitted by the core node 103 is set to zero. When a transmission interval of the route setting packet can be constant, random or on-demand in response to a command from the server on the wired backbone network 101.

At first, the relay nodes 104 to 106 check whether the route setting packet arrived or not (step S81 of Fig. 18). If the route setting packet is not detected, the process returns to step S81.

For detection of arrival of the route setting packet in the relay nodes 104 to 106, carrier sense or the like is used. When arrival of the route setting packet is detected (step S81 of Fig. 18), the relay nodes 104 to 106 makes reference to the receiver side relay node ID contained in the route setting packet for checking whether the receiver side relay node ID matches with the own node ID or not (step S89 of Fig. 18).

When the receiver side relay node ID matches with the own node ID, the relay nodes 104 to 106 record the ID of the node broadcasted the route setting packet, namely the sender node ID contained in the route setting packet in the relay node list (step S90 of Fig. 18).

The relay node list is a table indicating the down-link receiver side node number and is constructed as shown in Fig. 7. The relay node list is used as receiver side node list upon the down-link data packet relaying which was discussed befor.

Each receiver side relay node ID contained in the relay node list can be forgotten (erased) after elapse of a given period. For instance, when new relay node is added in the cell, when existing node moved or when new building is constructed in the cell, re-establishment of the relay route becomes necessary. In order to adapt to re-established relay route, each receiver side relay node ID contained in the relay node list may be intentionally forgotten after expiration of the given period.

If judgment is made that the relay node ID does not match with the own node ID, the relay nodes 104 to 106 takes a path loss as measured at the timing of judgment, as a path loss Ln (n is specific number of the sender node of the route setting packet) between the node broadcasted the route setting packet and the own node (step S82 of Fig. 18). Measurement of the path loss is generally performed upon reception of the packet irrespective of the content. For measurement of the path loss, reception power of the packet or the like may be used. In order to facilitate measurement of the path loss, transmission power of the route setting packet can be fixed. It should be noted that n represents a node number and the node number n is designated by the sender node ID contained in the route setting packet, as shown in Fig. 2.

The relay nodes 104 to 106 read the metric Mr, n contained in the received route setting packet (step S83 of Fig. 18).

Here, the update metric Mr,n is given as a sum of the path loss. The relay node 104 to 106 designate the update metric Mn from the path loss Ln and the metric Mr,n measured in step S82. Here, the update metric Mn is given by the sum of the trasmission loss Ln and the metric Mr,n. The relay nodes 104 to 106 store the update metric Mn calculated through the foregoing process (step S84 of Fig. 18).

It should be noted that, among stored update metric, the update metric maintained for a period in excess of a given period can be forgotten (erased). On the other hand, the stored metric is always the most recent one. Namely, when the update metric corresponding to the node n as the sender of the route setting packet was stored in the past, the old metric is re-written with the new update metric obtained at step S84.

The relay nodes 104 to 106 compares the update metric Mn with update metric corresponding to all route setting packet received in the past to determine the sender node m (m is number specific to the node) as the smallest metric (step S85 of Fig. 18). If the sender node m is the same as current up-link receiver side relay node of own node and $n \neq m$(step@ S86 of Fig. 18), process is returned to step S81 and transmission of the route setting packet is not newly performed.

If the sender node m having the minimum metric is not the same as the current up-link receiver side node or $n = m$

(step S86 of Fig. 18), the update metric Mn is set in the metric M contained in the metric A04, and the sender node m is registered as up-link receiver side relay node of own node (step S87 of Fig. 18). Namely, the route setting packet is transmitted when the up-link receiver side relay node is changed or when the route setting packet is received from the same up-link receiver side relay node even if up-link receiver side relay node is not changed.

The relay nodes 104 to 106 sets the metric M set as set forth above as metric and transmit the route setting packet with including the necessary information in respective of other items shown in Fig. 2 (step S88 of Fig. 18).

Upon reception of the route setting packet, the relay nodes 104 to 106 may return an achnowledgement signal for accuracy. Since the route setting packet is the control packet directed to unspecified node, the relay nodes 104 to 106 may receive the achnowledgement signals from a plurality of nodes after transmission of the route setting packet. When the relay nodes 104 to 106 does not receive returned acknowledgement signal at all, re-transmission of the route setting packet may be performed.

In Fig. 18, two kinds of metrics may be used as set forth above upon detection of the node having the minimum metric at step S85. Fig. 19 is a flowchart showing an example of another

process in a part of the relay route setting process in another embodiment of the present invention, in which steps S91 and S92 shown in Fig. 19 can be replacement of step S85 in Fig. 18. For example, in case of combinations of the first and second metrics are route A = (3 hops, 100 dB), route B (3 hops, 91 dB) and route C (4 hops, 85 dB) are present. At first, the route C having large hop count as the first metric is eliminated. Subsequently, by comparing the second metrics, the route B is selected.

As set forth above, the present invention can select the route having the minimum path loss in the entire relay route and can set the relay route resistive against interference by selecting the relay nodes to have total path loss in at least one of the routes between the relay nodes and between the relay node and the core node, in the radio network including the core node connected to the wired network, the relay node relaying at least one of the down-link data packet transmitted from the core node and the up-link data packet, and a terminal station capable of transmission and reception of the data packet with the core node and the relay node.

On the other hand, the present invention uses two kinds of metrics upon performing route control and calculates update metric with weighting to permit more flexible route setting to facilitate generation of the route having characteristics

expected by the network designer.

Although the present invention has been illustrated and described with respect to exemplary embodiment thereof, it should be understood by those skilled in the art that the foregoing and various other changes, omission and additions may be made therein and thereto, without departing from the spirit and scope of the present invention. Therefore, the present invention should not be understood as limited to the specific embodiment set out above but to include all possible embodiments which can be embodied within a scope encompassed and equivalent thereof with respect to the feature set out in the appended claims.